

Faster Runtime Verification during Testing via Feedback-Guided Selective Monitoring

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1. Background and Problem

Runtime Verification (RV)

- Monitors program executions against formal specifications (specs)
- Found hundreds of bugs regarding correct JDK API usage [1-3]

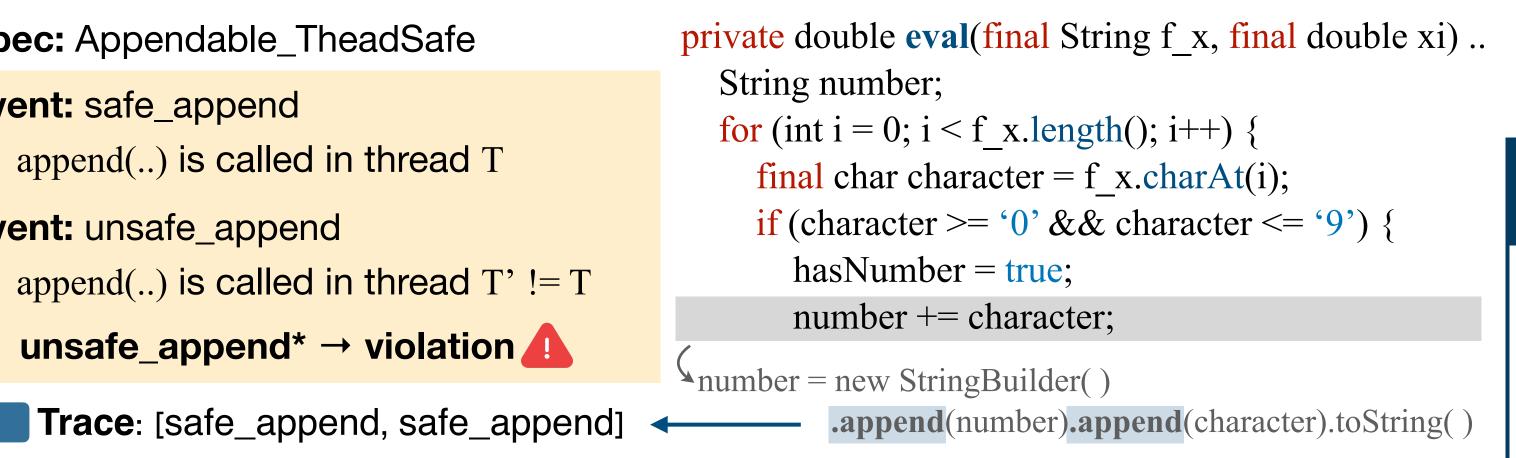
Problem: RV has runtime overhead which can be as high as 5,000x compared to unit testing only, or 27 hours [4]

- There exist two types of specs in our style of RV (MOP)
- Each type of specs incurs runtime overhead for different reasons

1. Parametric Specs

- RV creates a monitor for each set of spec-related objects
- 99.87% monitors are redundant for bug finding [4]. E.g.:

Spec: Appendable_TheadSafe **Event:** safe_append • append(..) is called in thread T **Event:** unsafe_append • append(..) is called in thread T' != T unsafe_append* → violation ____



RV creates 68,000,157 parametric monitors that check the same trace!

2. Our Approach: Valg

Yey Idea: Selective Monitoring!

- Use reinforcement learning for selective parametric monitor creation
- Use violation feedback to selectively signal non-parametric events

Selective Parametric Monitor Creation

Learning objective: Reduce redundant monitors and preserve unique ones

- 1. Formulation as a two-armed bandit problem
- **:** "Hyperparameters"
- Actions: {'create', 'ncreate'} // creating a monitor or not
- Binary reward for 'create' and continuous reward for 'ncreate'

 $R_{\text{create},t} \doteq 0 \text{ if (trace}_t \text{ is redundant) else } 1$ $R_{\text{ncreate},t} \doteq \frac{\sum_{k=0}^{t-1} 1(\text{trace}_k \text{ is redundant})}{\sum_{k=0}^{t-1} 1(\text{trace}_k \text{ is observed})}$

- 2. Selection based on action-value method Learning Rate Exploration Prob.
 - Estimate rewards using exponential-recency weighted average

 $Q_{n+1} \doteq Q_n + \alpha (R_n - Q_n)$, where α is a learning rate

• Enable stochastic exploration (vs. exploitation) using ϵ -greedy

 $A_t \leftarrow \begin{cases} \arg\max_a Q_t(a) & \text{with probability } 1 - \epsilon \\ \text{random action } a & \text{with probability } \epsilon \end{cases}$

- 3. Convergence logic for the learning Convergence Threshold
 - Heuristic: If the absolute difference in estimated values is close to 1
- 4. Initial value selection Initial Values
 - Optimistic value for 'create' to encourage monitor creation at early stages

@ iteration 1 @ iteration 2 **Trace**: [safe_append, safe_append] @ iteration 3 @ iteration 68,000,157

@ iteration 1 Trace: [safe_append, safe_append] [[create] 🔃 Trace: [safe_append, safe_append] 🛄 @ iteration 2 **Trace**: [safe_append, safe_append][€]√ [create] @ iteration 3 **Trace**: [safe_append, safe_append] [ncreate] **Trace**: [safe_append, safe_append] @ iteration 68,000,157

Trace: [safe_append, safe_append] [[ncreate] | Trace: [safe_append, safe_append] State-of-the-Art (SoTA)

Our Technique (Valg)

Valg reduces 68,000,157 created monitors to 2 monitors

Selective Non-Parametric Event Signaling

- Valg tracks the violation status of each event location
- Valg does not signal the event if a violation was already detected

Valg reduces 260,000,000 signaled events to 1 event

2. Non-Parametric Specs

- RV creates only one monitor for <u>all</u> spec-related objects or static calls
- Over 99.99% of signaled events are redundant for bug finding. E.g.:

Spec: Math_ContendedRandom

Event: onethread_use Math.random() is called in thread T

Event: otherthread use

• Math.random() is called in thread T' != T otherthread_use* → violation ____

public static String generateData(int byteSize) ...

StringBuilder b = new StringBuilder(byteSize *2); for (int i = 0; i < byteSize; i++) { if (Math.random() * 100 > 98) { // appends a terminating character to b

Trace: [onethread_use, onethread_use, ...]

RV signals 260,000,000 non-parametric events that check the same logic!

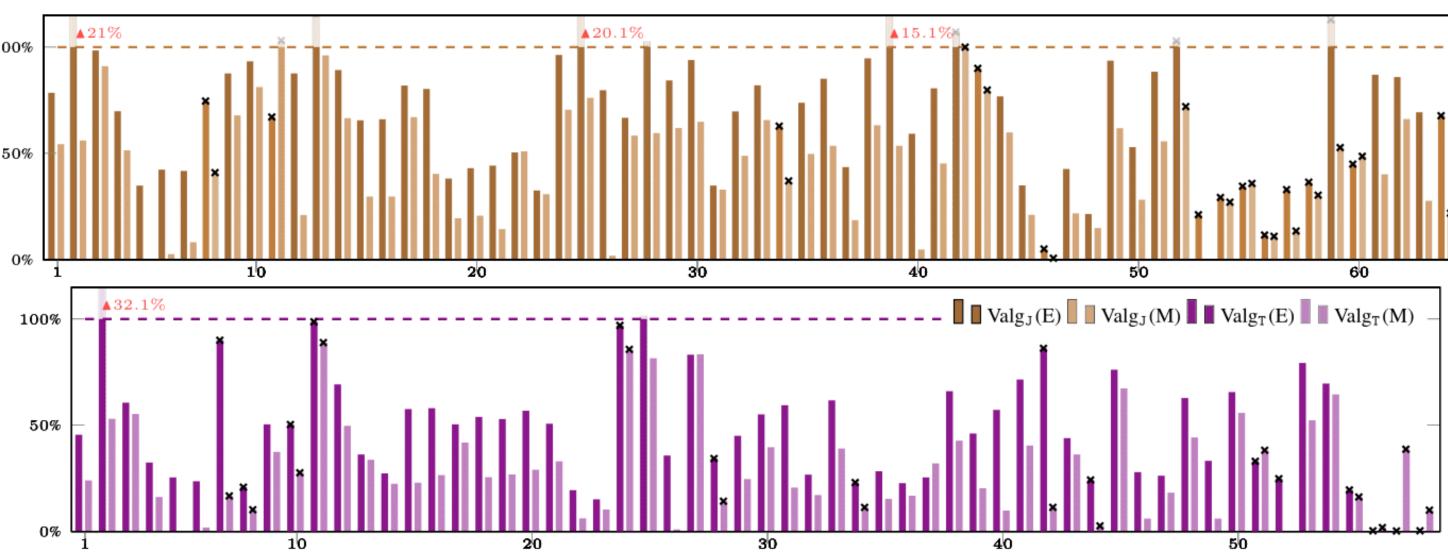
[2] Legunsen et al., "How Effective are Existing Java API Specifications for Finding Bugs during Runtime Verification?," JASE 2019 [3] Miranda et al., "Prioritizing Runtime Verification Violations," ICST 2020

[4] Guan and Legunsen, "An In-Depth Study of Runtime Verification Overheads During Software Testing," ISSTA 2024

3. Evaluation Results

Valg vs. SoTA Techniques (JavaMOP and TraceMOP [5])

Setup: 64 Java open-source projects, 160 JDK API specs



Compared to JavaMOP and TraceMOP,

Overhead. Valg is up to 20.2x (4.3 hrs) and 551.5x (24.3 hrs) faster Valg reduces 3.02 days down to 11.6 minutes for three projects

Violations. Valg preserves **99.6**% of the original violations

[5] Guan and Legunsen, "TraceMOP: An Explicit-Trace Runtime Verification Tool for Java," FSE Demo, 2025

RQ2 Valg vs. {10%, 50%} Random Sampling

• **Setup**: 20 Java open-source projects, 160 JDK API specs

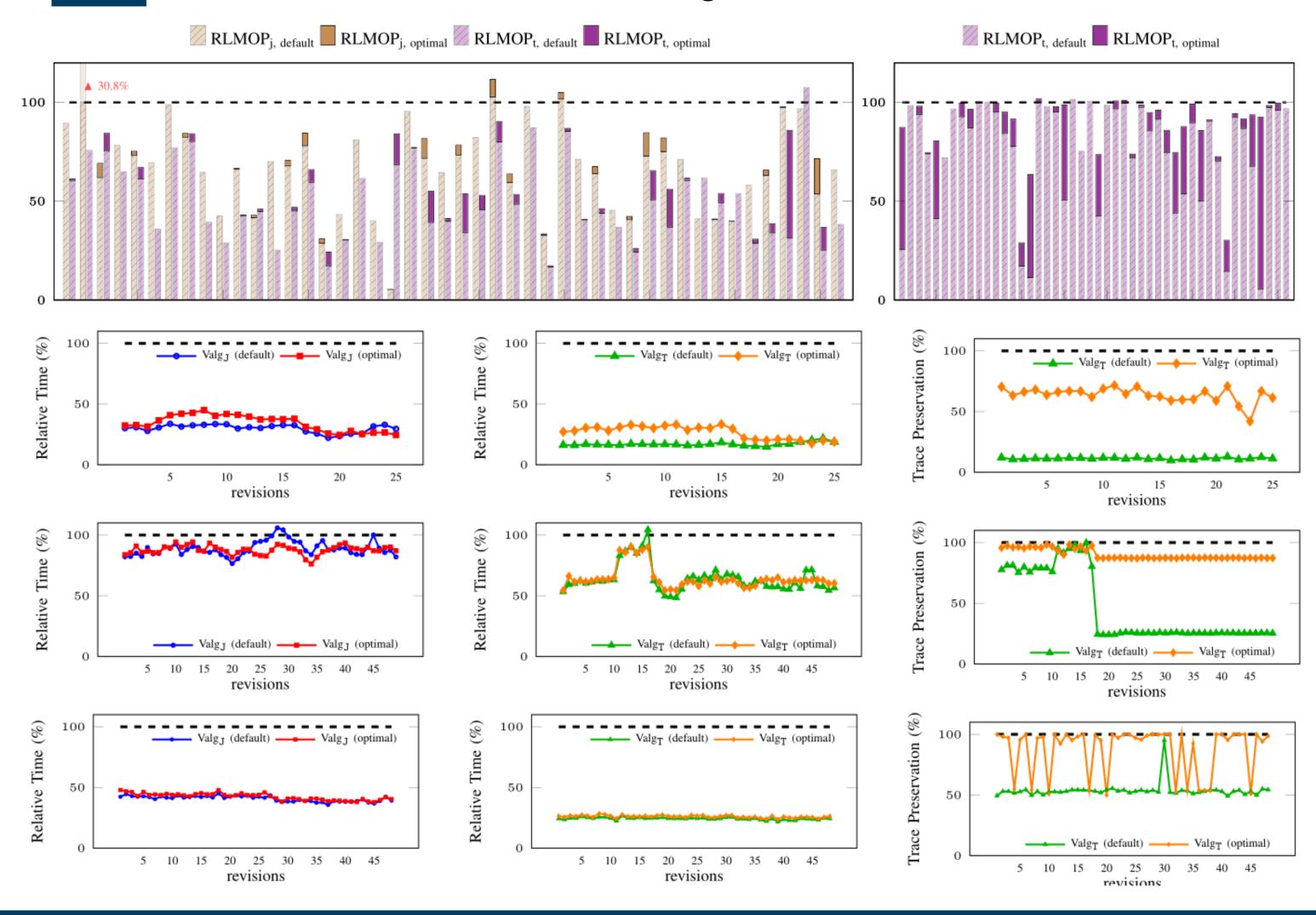
Compared to {10%, 50%} random sampling,

Overhead. Valg is up to 11.8x (4.3 mins) and 20.1x (22.0 mins) faster Violations. Valg preserves 66.1% and 14.4% more violations

RQ3 What is the impact of hyperparameter tuning?

Unique Traces. Valg's preservation ratio improves from 76.7% to 95.1%

RQ4 How effective and efficient is Valg as code evolves?



4. Discussion and Future Work

Discussion: Comparison with evolution-aware RV, ablation study, memory overhead Future Work: Valg opens up a new research direction for learning-based RV

Different algorithms, further study on hyperparameter tuning, hyperparameter learning